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# EEL 4783: Hardware/Software Co-design with FPGAs

## Lecture 11: Hardwar/Software Partitioning \*

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\* Adopted from G. Khan COE718:

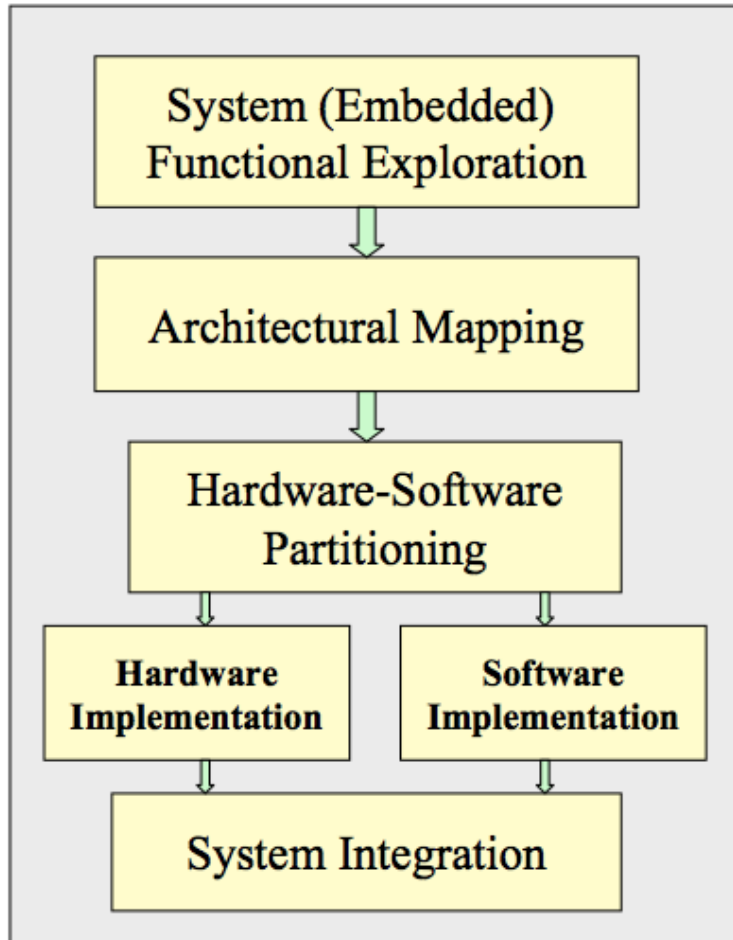
# Embedded Computing System

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- HW and SW used to be designed separately
- Co-Design an increasingly important object, because
  - Performance
  - Cost
  - Complexity
  - ?

# HW/SW Co-Design

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- Functional exploration: Define a desired product's requirements and produce a specification of the system behavior.
- Map this specification
- Partition the functions between silicon and code, and map them
- Integrate system

# HW/SW Co-Design

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- Co-Specification: Describe system functionality at the abstract level
- System description is converted into a task graph representation
- HW-SW Partitioning: Take the task graph and decide which components are implemented where/how ?
  - i.e. Dedicated hardware Software

# HW/SW Co-Design

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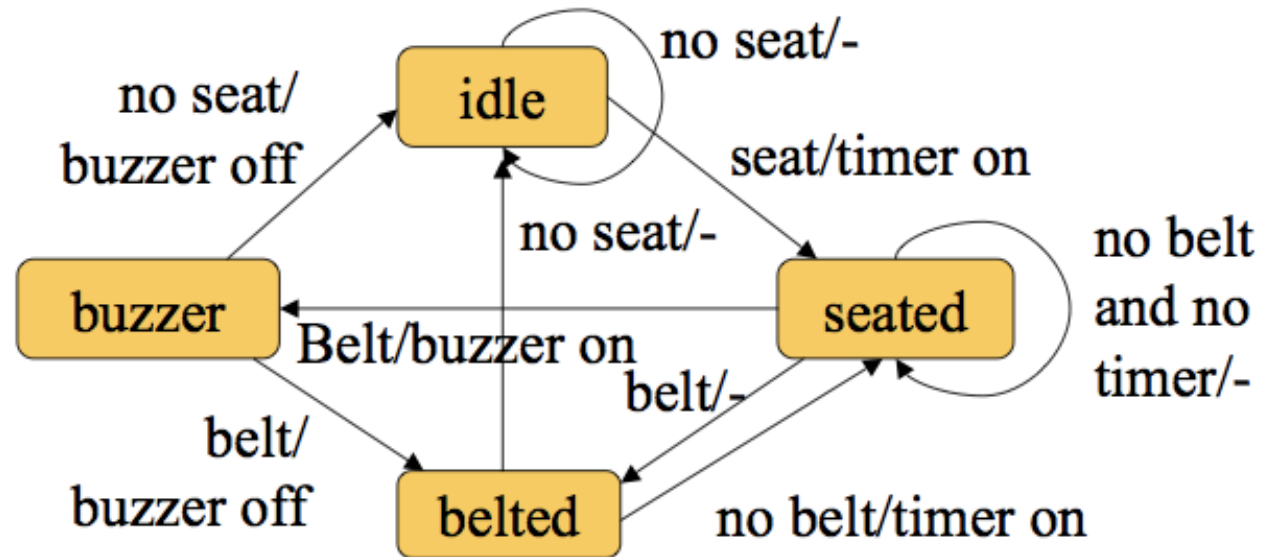
- Both textual and graphical representation like DAG are used to describe system.
- Analyzes task graph to determine each task's placement
- Many partitioning algorithms being developed
- Major problem involves the computation time of the algorithm.

# System Design Patterns

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- Design Pattern: A generalized description of the design of a certain type of program that can also be used for system representation and hardware-software partitioning.
- State Diagram  
Data Flow Graph  
Control Data Flow Graph (CDFG)
  - Others
    - Directed Acyclic Graph (DAG) similar to DFG
    - Directed Acyclic Data Dependence Graph with Precedence (DADGP)

# State Machine: Seat-belt System



```
switch (state) {  
    case IDLE: if (seat) { state = SEATED; timer_on = TRUE; } break;  
    case SEATED: if (belt) state = BELTED;  
                 else if (timer) state = BUZZER; break;  
    .....  
}
```

# Data Flow Graph (DFG)

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- DFG does not represent control
- It models the Basic Block: code or a system block with one entry and exit
- Describes the minimal ordering requirements on operations



# Data Flow Graph: Software Module

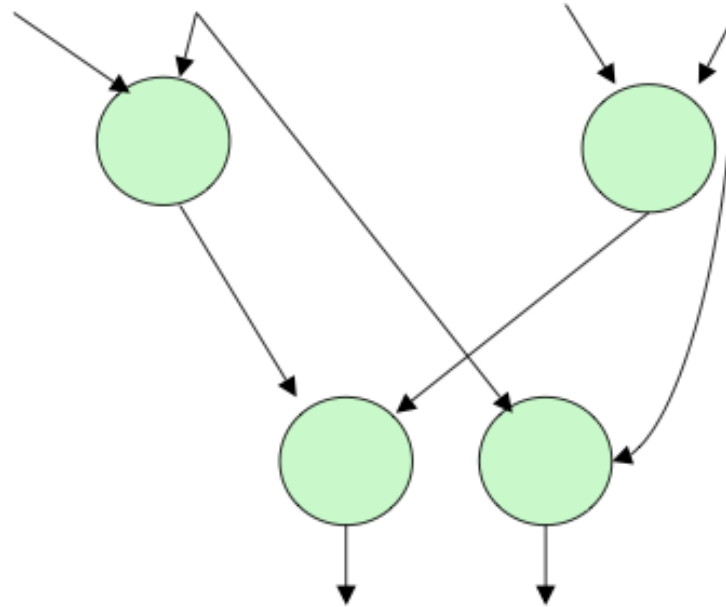
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$x = a + b;$

$y = c - d;$

$z = x * y;$

$y1 = b + d;$



**DFG**

# Control Data Flow Graph

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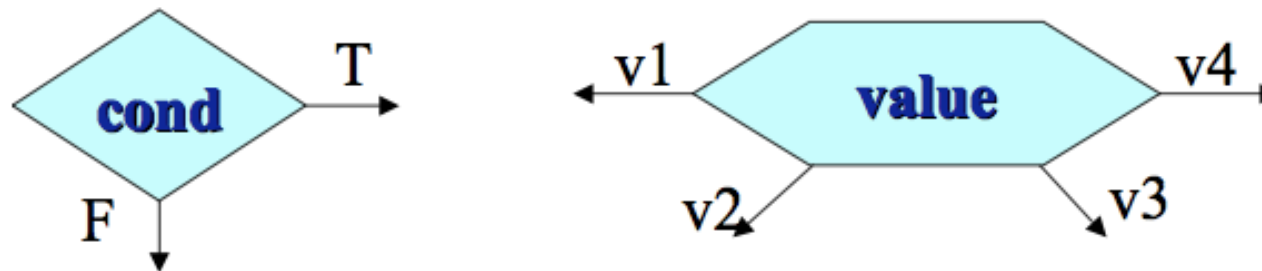
**CDFG:** represents control and data.

- Uses data flow graphs as components.
- Two types of nodes:

- **Data Flow Node encapsulate a DFG**

```
x = a + b;  
y = c + d
```

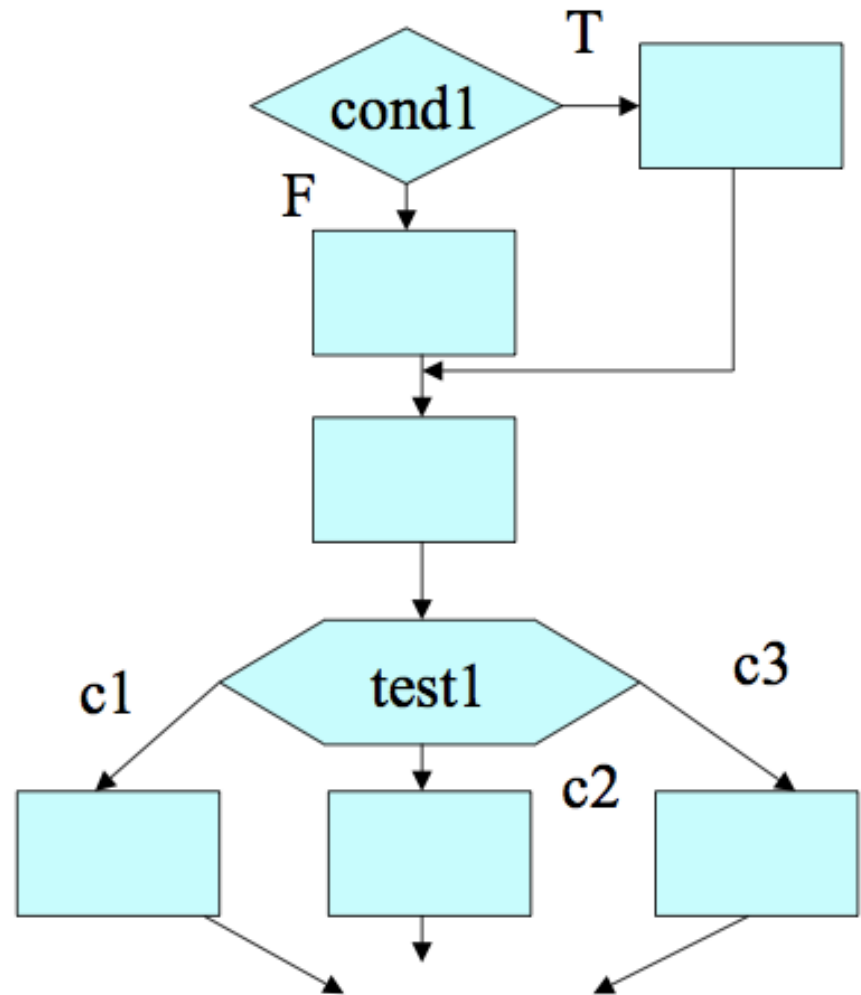
- **Decision Nodes**



**Equivalent Forms**

# Control Data Flow Graph Example

```
if (cond1) bb1();  
else bb2();  
bb3();  
switch (test1) {  
  case c1: bb4(); break;  
  case c2: bb5(); break;  
  case c3: bb6(); break;  
}
```



# Scheduling and Partitioning

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- The main input to scheduling for partitioning is a graph representation in the form of DFG.
- Complex designs contain thousands of both control and data processing operations ranging from:
  - Complex arithmetic operations or logic-level bit- operations.
  - All the above interleaved operations by multiple control operations and loops.
- Such designs contain thousands of data-dependencies, basic blocks and control paths.

# DFG-based Scheduling & Partitioning

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- Data-flow based scheduling techniques extract parallelism from the input description
- Schedule operations in parallel to satisfy the constraints.
- Two most common DF-based scheduling methods.
  - List Scheduling (LS): Minimize the number of control steps under resource constraints.
  - Force-directed Scheduling (FDS): Minimize the number of resource constraints under a fixed number of control steps.

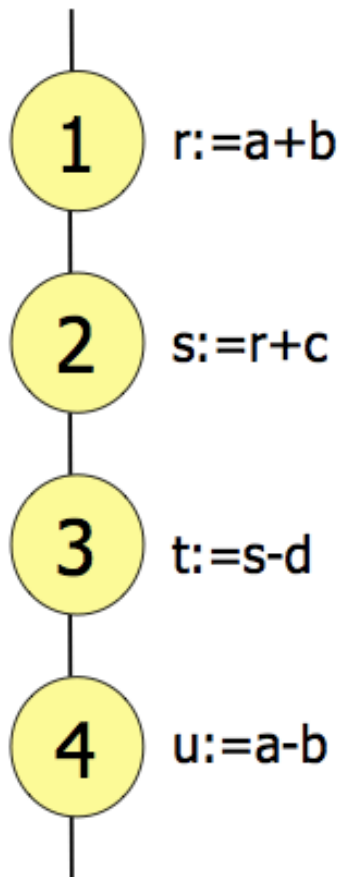
# Data Flow: DF-Scheduling

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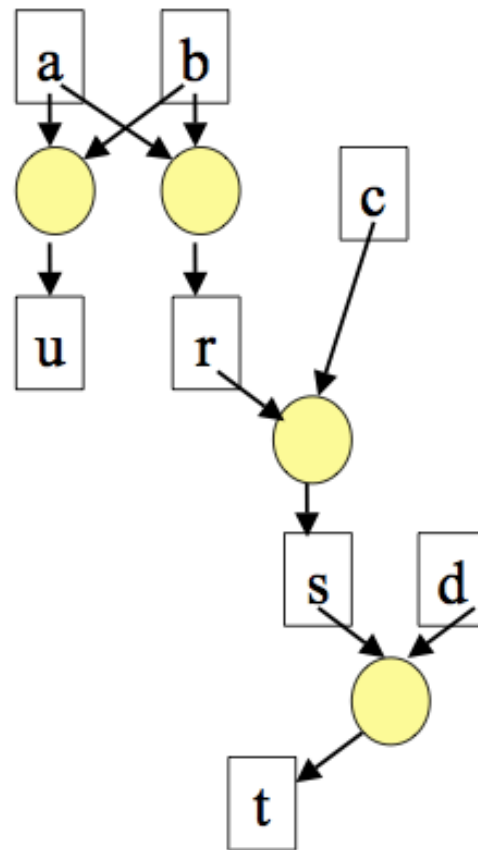
- List scheduling algorithm uses a cost function to select the operation to be scheduled from a list.
- DF-approach provides flexible cost function and it can be easily adapted to generate resource-constraint as well as time-constraint schedules.
- The cost function can represent any design measure such as HW area, delay, etc.
- The result is only as good as the cost function.
- DF-based algorithms can analyze all the parallelism in the DFG independently.

# DF- Scheduling Example

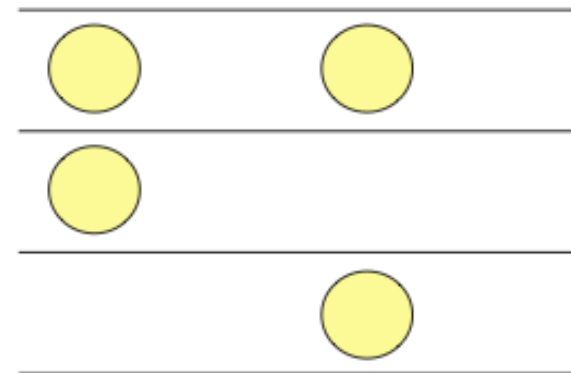
CFG



DFG



DFG-Schedule



# Control Flow: CF-Scheduling

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- Analyze the sequences of operations in CFG called control flow paths and schedule the CFG with minimum number of control steps in each path.
- Path-based scheduling is one of the main example of this scheme.
- Analyze all the paths in the CFG and schedule each of them independently.
- It minimizes the number of control steps in each path rather than minimizing the number of states.
- Paths in CFG come from loops and conditional operations.



# Partitioning Approaches

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- Simple one CPU and single ASIC architecture is the most common.
- Early approaches Initially assume all tasks mapped to software
- Move tasks to HW incrementally until system requirements are met.
- Other early approaches: Initially all tasks are mapped to dedicated hardware.
- Move tasks incrementally to SW until system requirements are met.

# Optimal Partitioning

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- Exhaustive approaches are characterized by attempting all possible combinations there by always selecting the best option.
- Exhaustive approaches are generally computationally intensive, consume huge-time in the range of hours or even days to find an optimal partition.
- Limited to small task graphs
  - Large telecom or other embedded systems can have upto 4000 nodes

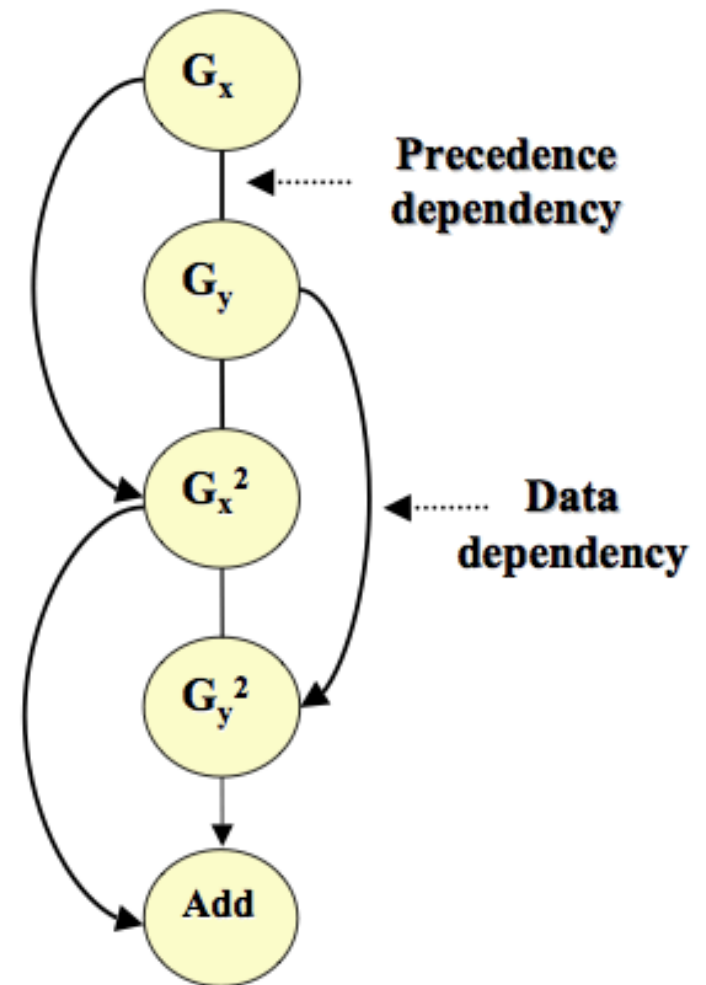
# Edge Detection Example

Pair of masks are convolved to estimate gradients,  $G_x$  and  $G_y$

$$\text{Overall } G^2 = (G_x^2 + G_y^2)$$

## HW-SW Library

Operation	SW EXE (ms)	HW EXE (ms)	HW Area (gates)
Gradient ( $G_x$ or $G_y$ )	9.4	1.4	1200
Square	5.2	0.9	500
Add	3.88	0.3	100



# SOBEL Edge Detection

## SOBEL masks

-1	0	+1
-2	0	+2
-1	0	+1

Gx

+1	+2	+1
0	0	0
-1	-2	-1

Gy

Input Image

a <sub>11</sub>	a <sub>12</sub>	a <sub>13</sub>		
a <sub>21</sub>	a <sub>22</sub>	a <sub>23</sub>		
a <sub>31</sub>	a <sub>32</sub>	a <sub>33</sub>		

Mask

m <sub>11</sub>	m <sub>12</sub>	m <sub>13</sub>
m <sub>21</sub>	m <sub>22</sub>	m <sub>23</sub>
m <sub>31</sub>	m <sub>32</sub>	m <sub>33</sub>

Output Image

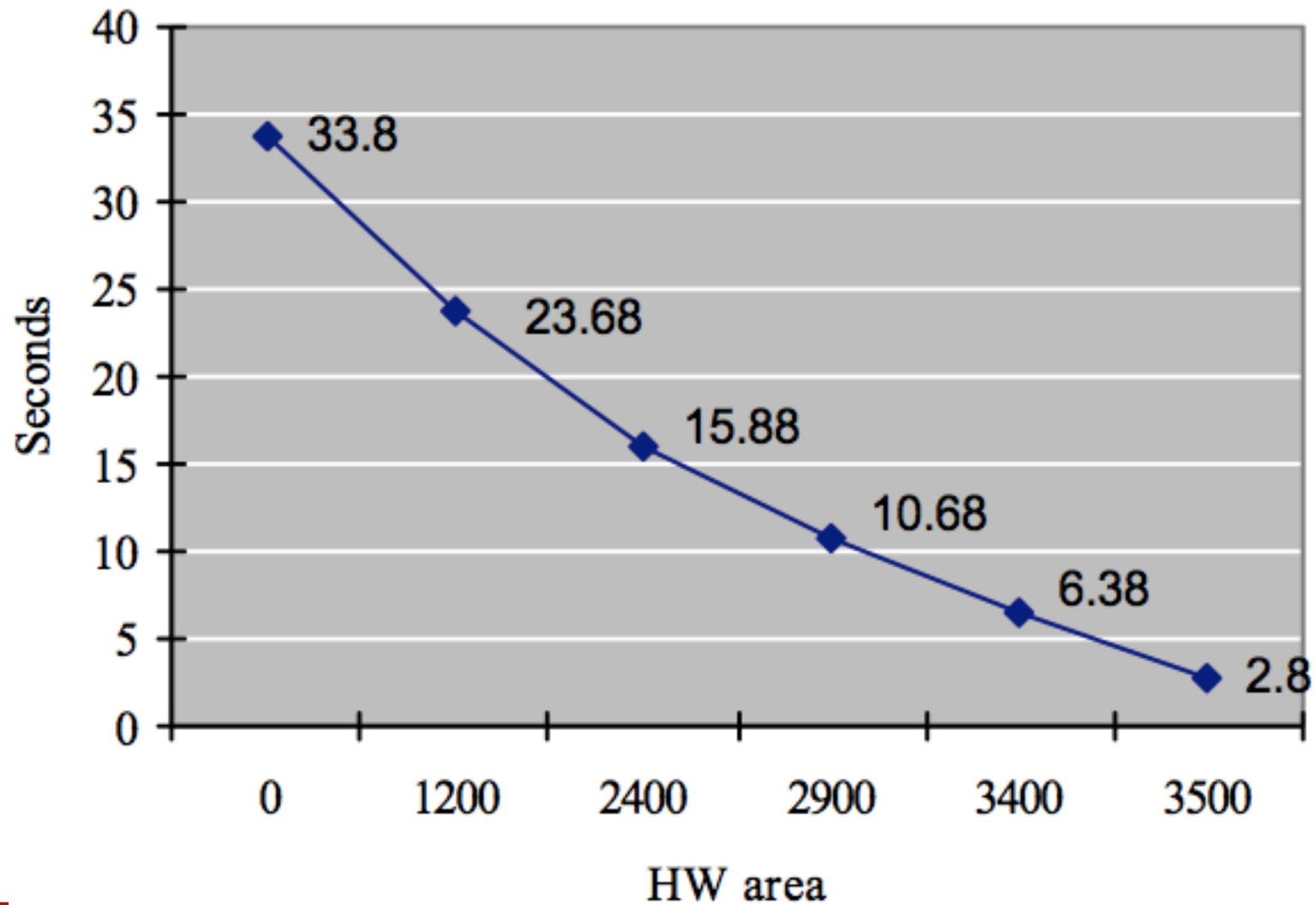
b <sub>11</sub>	b <sub>12</sub>	b <sub>13</sub>		
b <sub>21</sub>	b <sub>22</sub>	b <sub>23</sub>		
b <sub>31</sub>	b <sub>32</sub>	b <sub>33</sub>		

$$b_{22} = (a_{11} * m_{11}) + (a_{12} * m_{12}) + (a_{13} * m_{13}) + (a_{21} * m_{21}) + (a_{22} * m_{22}) + (a_{23} * m_{23}) + (a_{31} * m_{31}) + (a_{32} * m_{32}) + (a_{33} * m_{33})$$

# SOBEL Edge Detection

```
main() {  
    unsigned char image_in[ROWS][COLS];  
    unsigned char image_out[ROWS][COLS];  
    int r, c; /* row and column array counters */  
    int pixel; /* temporary value of pixel */  
        /*filter the image and store result in output array */  
    for (r=1; r<ROWS-1; r++)  
        for (c=1; c<COLS-1; c++) { /* Apply Sobel operator. */  
            pixel = image_in[r-1][c+1]-image_in[r-1][c-1]  
                + 2*image_in[r][c+1] - 2*image_in[r][c-1]  
                + image_in[r+1][c+1] - image_in[r+1][c-1];  
            /* Normalize and take absolute value */  
            pixel = abs(pixel/4);  
            /* Check magnitude */  
            if (pixel > Threshold)  
                pixel= 255; /*EDGE_VALUE;*/  
            /* Store in output array */  
            image_out[r][c] = (unsigned char) pixel;  
        }  
    }  
}
```

# SOBEL Edge Detection



# Final issues

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- Come by my office hours (right after class)
- Any questions or concerns?